

# Self-stabilizing Systems in Spite of Distributed Control

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# Overview

- Motivation
- Complications
- Concepts
- Solutions

# Motivation

- Parallel, distributed, multi-threaded computing etc. needs synchronization
- How to guarantee the correctness?
- How to recover from the incorrect system?

# Motivation

- Correctness
- The system is in a legitimate state
- Recovery
- The system will reach a legitimate state

# Classification

- System state is saved in common place
- No common store for synchronization

# Concepts

- Neighbors
- Distributed Control
- Privileges
- Self-stabilizing

# Neighbors

- Processes with direct communication capability
- Worse: a small subset of processes

# Distributed Control

- No common store for mutually exclusive access
- The system state is saved in variables across various processes



# Privileges

- For machines
- Boolean function of states (L,S,R)
- True -> Present -> Move -> New state
- Nondeterministic

# Legitimate System from Privilege

- One or more privileges are present
- Legitimate state moves to another legitimate state
- Each privilege must present in at least one legitimate state
- There is a path between any legitimate states

# Self-stabilizing

- Regardless of initial state and privilege selected
- At least one privilege present
- The system moves into a legitimate state after a finite number of steps

# Solutions

- Assumptions:  $N+1$  machines in a ring
- Nr. 0 is the bottom machine
- Nr.  $N$  is the top machine
- Solutions:

K-state machines

Four-state machines

Three-state machines

# K-state Machines

- For the bottom machine:

If  $L=S$  then  $S := (S+1) \bmod K$  fi

- For others

If  $L \neq S$  then  $S := L$  fi

Where  $L, R$  are states of left, right neighbor

And  $0 \leq S < K$

# Four-state machines

- Each state has xS and upS
- For the bottom machine:  
upS=true
- For the top machine:  
upS=false

# Four-state machines

- The bottom machine

if  $x_S = x_R$  and non  $up_R$  then  $x_S := \text{non } x_S$  fi

- The top machine

If  $x_S \neq x_L$  then  $x_S := \text{non } x_S$  fi

- Others

If  $x_S \neq x_L$  then  $x_S := \text{non } x_S$ ;  $up_S := \text{true}$  fi;

If  $x_S = x_R$  and  $up_S$  and non  $up_R$  then  $up_S := \text{false}$  fi

# Three-state Machines

- The bottom machine

If  $(S+1) \bmod 3 = R$  then  $S := (S-1) \bmod 3$  fi

- The top machine

If  $L=R$  and  $(L+1) \bmod 3 \neq S$  then  $S := (L+1) \bmod 3$  fi

- Others

If  $(S+1) \bmod 3 = L$  then  $S := L$  fi;

If  $(S+1) \bmod 3 = R$  then  $S := R$  fi



# Question 1

- Are all systems “self-stabilizing”? Are all fault-tolerant systems “self-stabilizing”?

# Question 2

- The daemon selects privilege(s) for next move, so in real implementation, is the selection arbitrary? Or otherwise?