5

Control and Implementation of State Space Search

- 5.0 Introduction
- 5.1 Recursion-Based Search
- 5.2 Pattern-directed Search
- 5.3 **Production Systems**

- 5.4 The Blackboard Architecture for Problem Solving
- 5.5 Epilogue and References
- 5.6 Exercises

- Compare the recursive and iterative implementations of the depth-first search algorithm
- Learn about pattern-directed search as a basis for production systems
- Learn the basics of production systems
- The agent model: Has a problem, searches for a solution, has different ways to model the search

Summary of previous chapters

- Representation of a problem solution as a path from a start state to a goal
- Systematic search of alternative paths
- Backtracking from failures
- Explicit records of states under consideration
 - open list: untried states
 - closed lists: to implement loop detection
- open list is a *stack* for DFS, a *queue* for BFS

Function depthsearch algorithm

function depthsearch;

% open & closed global

begin if open is empty then return FAIL; current_state := the first element of open; if current_state is a goal state then return SUCCESS else begin open := the tail of open; closed := closed with current_state added; for each child of current state % build stack if not on closed or open then add the child to the front of open end; % recur depthsearch end.

- for clarity, compactness, and simplicity
- call the algorithm recursively for each child
- the open list is not needed anymore, activation records take care of this
- still use the closed list for loop detection

Function depthsearch (current_state) algorithm

function depthsearch (current_state);

begin if current_state is a goal then return SUCCESS; add current_state to closed; while current_state has unexamined children begin child := next unexamined child; if child not member of closed then if depthsearch(child) = SUCCESS then return SUCCESS end; return FAIL end % closed is global

% search exhausted

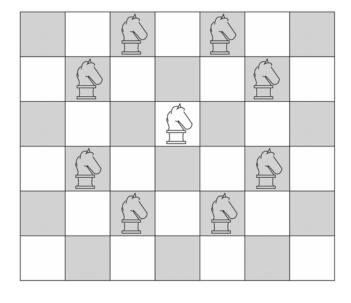
- \bullet use modus ponens on rules such as $q(X) \rightarrow p(X)$
- if p(a) is the original goal, after unification on the above rule, the new subgoal is q(a)

function pattern_search (current_goal);

begin if current_goal is a member of closed % test for loops then return FAIL else add current_goal to closed; while there remain in data base unifying facts or rules do begin case current_goal unifies with a fact: return SUCCESS; current_goal is a conjunction $(p \land ...)$: begin for each conjunct do call pattern_search on conjunct; if pattern_search succeeds for all conjuncts then return SUCCESS else return FAIL end; current_goal unifies with rule conclusion (p in $q \rightarrow p$): begin apply goal unifying substitutions to premise (q); call pattern_search on premise; if pattern_search succeeds then return SUCCESS else return FAIL end; % end case end; end: return FAIL end.

A chess knight's tour problem

Legal moves of a knight



1	2	3
4	5	6
7	8	9

- move(1,8)move(6,1)move(1,6)move(6,7)move(2,9)move(7,2)move(2,7)move(7,6)
- move(3,4) move(8,3)
- move(3,8) move(8,1)
- move(4,9) move(9,2)
- move(4,3) move(9,4)

Move rules

Examples

- Is there a move from 1 to 8? Pattern_search(move(1,8)) success
- Where can the knight move from 2? Pattern_search(move(2,X)) {7/X}, {9/X}
- Can the knight move from 2 to 3? Pattern_search(move(2,3)) fail
- Where can the knight move from 5? Pattern_search(move(5,X)) fail

2 step moves

- ∀X,Y [path2(X,Y) ←∃Z [move(X,Z) ∧ move(Z,Y)]]
- path2(1,3)?
- path2(2,Y)?

3 step moves

- ∀X,Y [path3(X,Y) ←∃Z,W [move(X,Z) ∧ move(Z,W) ∧ move(W,Y]]
- path3(1,2)?
- path3(1,X)?
- path3(X,Y)?

General recursive rules

- ∀X,Y [path(X,Y) ←∃Z [move(X,Z) ∧ path(Z,Y)]]
- ∀X path(X,X)

- if the current goal is negated call pattern_search with the goal and return success if the call returns failure
- if the current goal is a conjunction call pattern_search for all the conjuncts
- if the current goal is a disjunction call pattern_search for all the disjuncts until one returns success

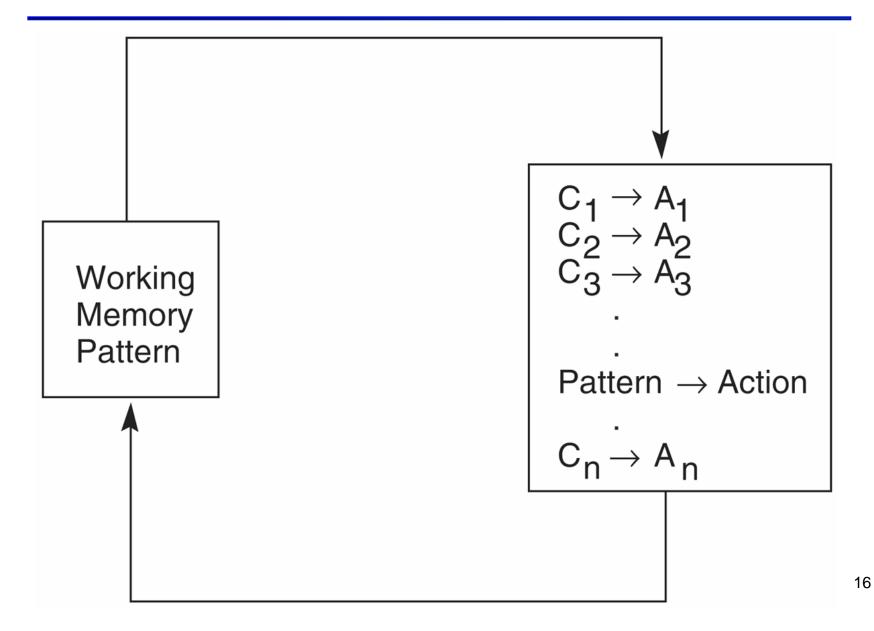
A *production system* is defined by:

• A set of *production rules* (aka *productions*): condition-action pairs.

• Working memory: the current state of the world

•The recognize-act cycle: the control structure for a production system Initialize working memory Match patterns to get the conflict set (enabled rules) Select a rule from the conflict set (conflict resolution) Fire the rule

A production system



Trace of a simple production system

Production set:

- 1. ba \rightarrow ab
- 2. ca \rightarrow ac
- 3. cb \rightarrow bc

Iteration #	Working memory	Conflict set	Rule fired
0	cbaca	1, 2, 3	1
1	cabca	2	2
2	acbca	2, 3	2
3	acbac	1, 3	1
4	acabc	2	2
5	aacbc	3	3
6	aabcc	Ø	Halt

17

The 8-puzzle as a production system

Start state: Goal state: 2 8 З 1 2 З 1 6 8 4 4 7 7 5 6 5

Production set:

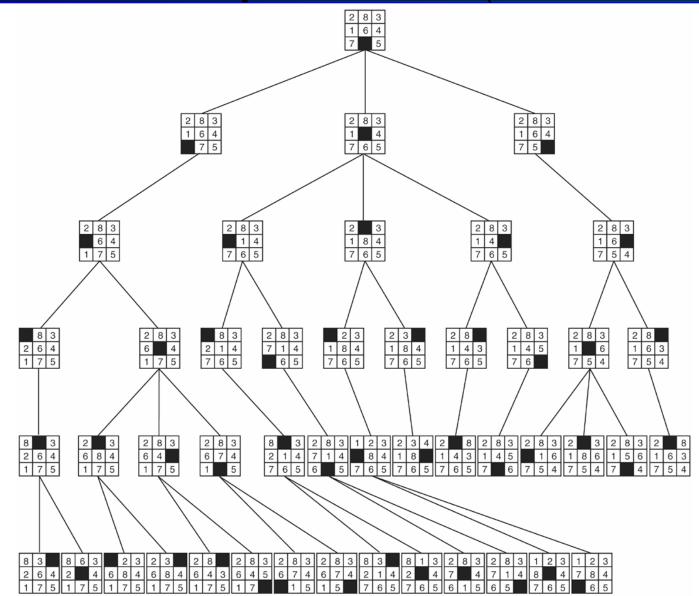
Condition	Action		
goal state in working memory blank is not on the left edge blank is not on the top edge blank is not on the right edge blank is not on the bottomedge			

Working memory is the present board state and goal state.

Control regime:

- 1. Try each production in order.
- 2. Do not allow loops.
- 3. Stop when goal is found.

Production system search with loop detection & depth bound 5 (Nilsson, 1971)

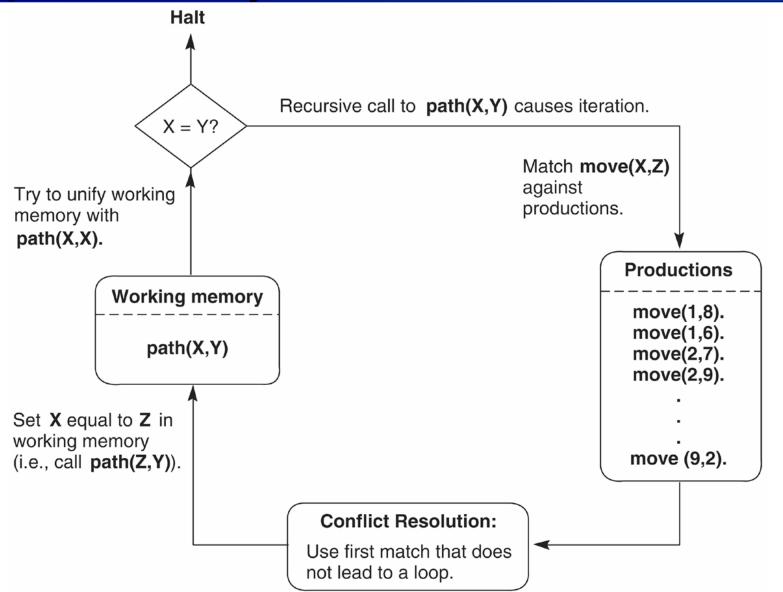


19

A production system solution to the 3×3 knight's tour problem

Iteration #	Working memory		Conflict set (rule #'s)	Fire rule
	Current square	Goal square		
0	1	2	1, 2	1
1	8	2	13, 14	13
2	3	2	5, 6	5
3	4	2	7, 8	7
4	9	2	15, 16	15
5	2	2		Halt

The recursive path algorithm: a production system



21

Data-driven search in a production system

Production set:

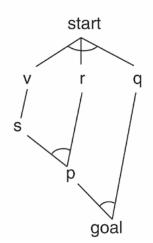
1. $p \land q \rightarrow goal$ 2. $r \land s \rightarrow p$ 3. $w \land r \rightarrow q$ 4. $t \land u \rightarrow q$ 5. $v \rightarrow s$

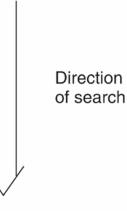
6. start $\rightarrow v \wedge r \wedge q$

Trace of execution:

Iteration #	Working memory	Conflict set	Rule fired
0	start	6	6
1	start, v, r, q	6, 5	5
2	start, v, r, q, s	6, 5, 2	2
3	start, v, r, q, s, p	6, 5, 2, 1	1
4	start, v, r, q, s, p, goal	6, 5, 2, 1	halt

Space searched by execution:





Goal-driven search in a production system

Production set:

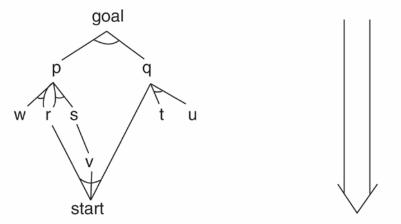
1. $p \land q \rightarrow goal$ 2. $r \land s \rightarrow p$ 3. $w \land r \rightarrow p$ 4. $t \land u \rightarrow q$ 5. $v \rightarrow s$

6. start $\rightarrow v \wedge r \wedge q$

Trace of execution:

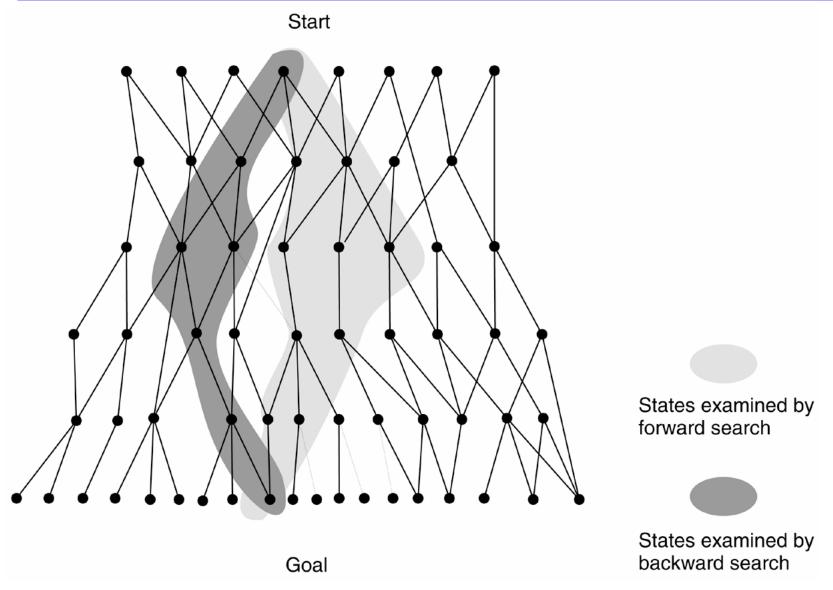
Iteration #	Working memory	Conflict set	Rule fired
0	goal	1	1
1	goal, p, q	1, 2, 3, 4	2
2	goal, p, q, r, s	1, 2, 3, 4, 5	3
3	goal, p, q, r, s, w	1, 2, 3, 4, 5	4
4	goal, p, q, r, s, w, t, u	1, 2, 3, 4, 5	5
5	goal, p, q, r, s, w, t, u, v	1, 2, 3, 4, 5, 6	6
6	goal, p, q, r, s, w, t, u, v, start	1, 2, 3, 4, 5, 6	halt

Space searched by execution:



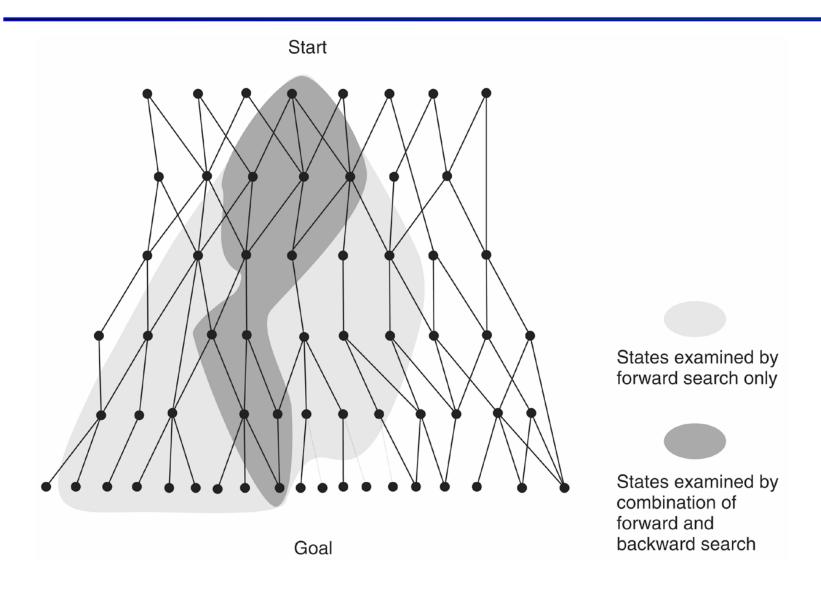
Direction of search

Bidirectional search misses in both directions: excessive search



24

Bidirectional search meets in the middle



Advantages of production systems

Separation of knowledge and control

A natural mapping onto state space search

Modularity of production rules

Pattern-directed control

Opportunities for heuristic control of search

Tracing and explanation

Language independence

A plausible model of human problem solving

Given a start state and a goal state

• State space search keeps the "current state" in a "node". Children of a node are all the possible ways an operator can be applied to a node

•Pattern-directed search keeps the all the states (start, goal, and current) as logic expressions. Children of a node are all the possible ways of using modus ponens.

• Production systems keep the "current state" in "working memory." Children of the current state are the results of all applicable productions.

Variations on a search theme

- **Bidirectional search**: Start from both ends, check for intersection (Sec. 5.3.3).
- **Depth-first with iterative deepening**: implement depth first search using a *depth-bound*. Iteratively increase this bound (Sec. 3.2.4).
- Beam search: keep only the "best" states in OPEN in an attempt to control the space requirements (Sec. 4.4).
- Branch and bound search: Generate paths one at a time, use the best cost as a "bound" on future paths, i.e., do not pursue a path if its cost exceeds the best cost so far (Sec. 3.1.2).