High Performance Dense Linear System Solver with Soft Error Resilience

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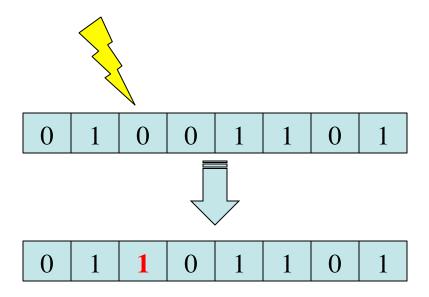


Agenda

- Soft error threat to the dense linear solver
 - LU factorization
 - Error propagation
- Error modeling
- Fault tolerant algorithm
- Performance Evaluation

Soft error

- Silent error due to radiation
 - Alpha particle
 - High energy neutron
 - Thermal neutron



- Outbreaks
 - Commercial computing system from Sun Microsystem in 2000
 - ASC Q supercomputer at Los Alamos National Lab in 2003

LU based linear solver

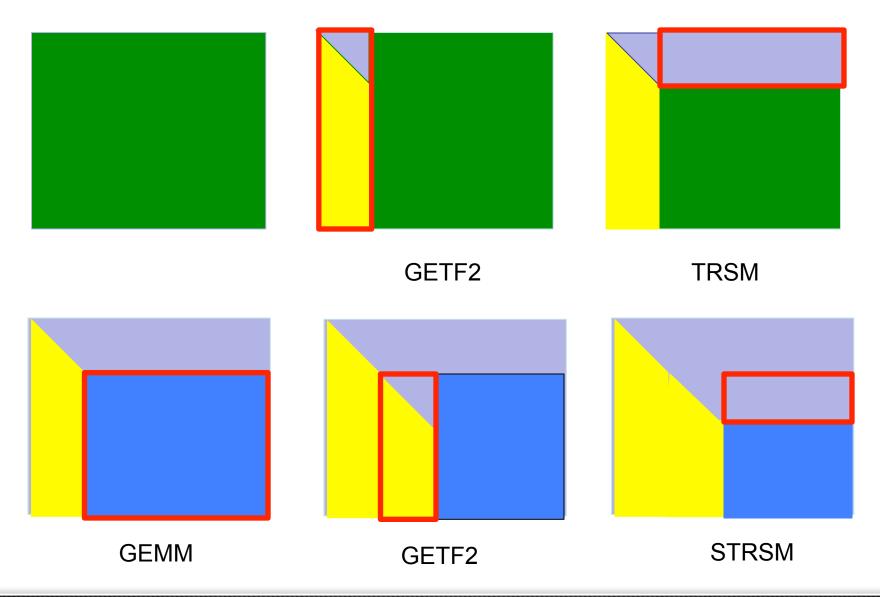
$$Ax = b$$

$$A = LU$$

$$x = U \setminus (L \setminus b)$$

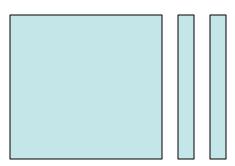
```
>> A=rand(4,4)
    0.6557
              0.6787
                        0.6555
                                   0.2769
    0.0357
              0.7577
                        0.1712
                                   0.0462
    0.8491
              0.7431
                        0.7060
                                   0.0971
    0.9340
              0.3922
                        0.0318
                                   0.8235
>> b=rand(4,1)
b =
    0.6948
    0.3171
    0.9502
    0.0344
>> [L,U]=lu(A)
                        0.9188
                                   1.0000
    0.7021
              0.5431
    0.0382
              1.0000
    0.9091
              0.5204
                        1.0000
    1.0000
U =
    0.9340
                        0.0318
              0.3922
                                   0.8235
              0.7427
                        0.1700
                                   0.0147
                        0.5886
                                  -0.6591
                                   0.2964
>> x=U\(L\b)
    0.4643
    0.3126
    0.5486
   -0.6549
>> norm(A*x-b)
ans =
   1.5701e-16
```

Block LU factorization



General work flow

- (1) Generate checksum for the input matrix as additional columns
- (2) Perform **LU factorization** WITH the additional checksum columns
- (3) <u>Solve Ax=b</u> using LU from the factorization (even if soft error occurs during LU factorization)
- (4) **Check** for soft error
- (5) **Correct** solution *x*

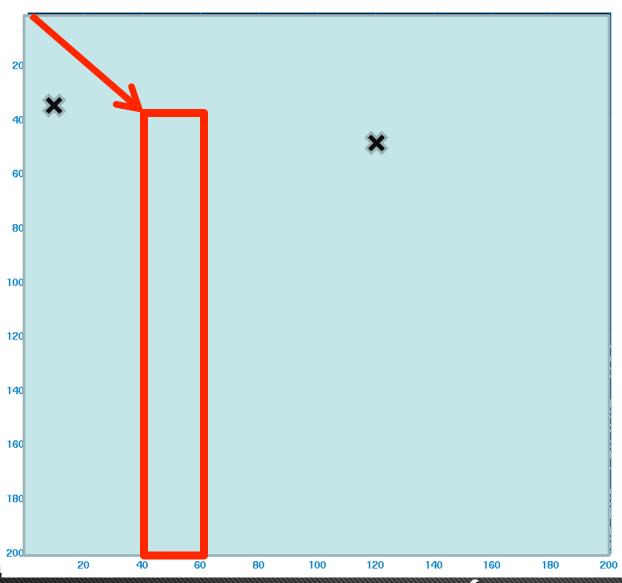


Why is soft error hard to handle?

Soft error occurs silently

Propagation

Example: Error propagation



Error location (using matlab notation and 1-based index)

Error strikes right before panel factorization of (41:200, 41:60),

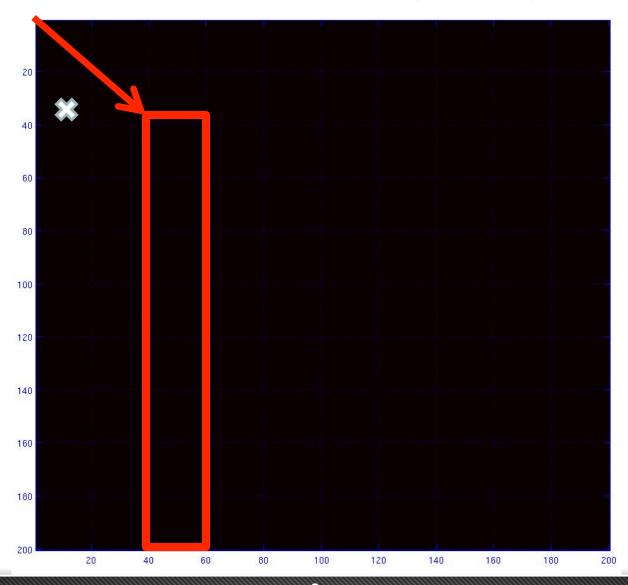
Case 1: Error at (35,10), in L area

Case 2: Error at (50,120), in A' area

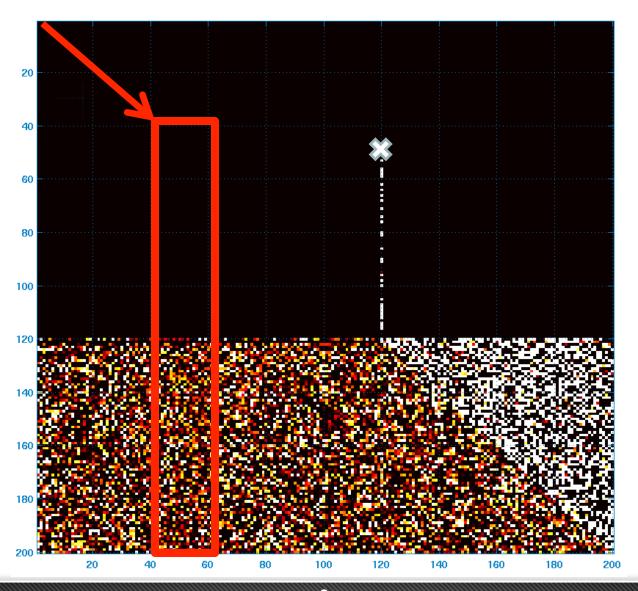
Note: Pivoting on the left of panel factorization is delayed to the end of error detection and recovery so that error in L area does not move



Case 1: Non-propagating error



Case 2: Propagating error



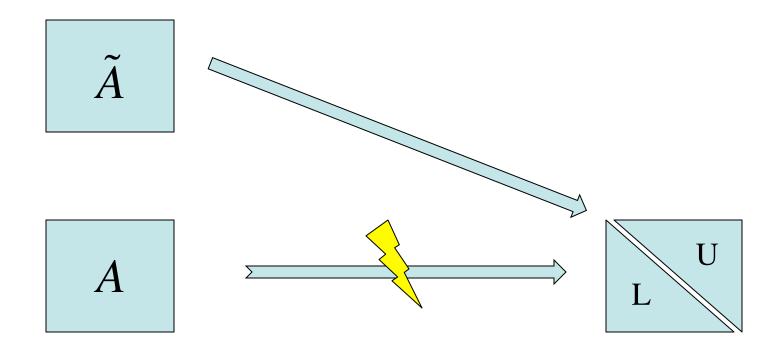
Soft error challenge

When?

Where?

Error modeling (for propagating error)

- When?
 - Answer: Doesn't really matter



Error modeling (for "where")

Input matrix

 \boldsymbol{A}

One step of LU
$$A_t = L_{t-1} P_{t-1} A_{t-1}$$

If no soft error occurs
$$U = (L_n P_n) \cdots (L_1 P_1) (L_0 P)_0 A_0$$

If soft error occurs at step t
$$\tilde{A}_t = L_{t-1}P_{t-1}A_{t-1} - \lambda e_i e_j^T$$

$$= L_{t-1}P_{t-1}(L_{t-2}P_{t-2}\cdots L_0P_0)A_0 - \lambda e_i e_j^T$$

Define an initial erroneous initial matrix \tilde{A}

$$\tilde{A} \cong (L_{t-1}P_{t-1}L_{t-2}P_{t-2}\cdots L_{0}P_{0})^{-1}\tilde{A}_{t}$$

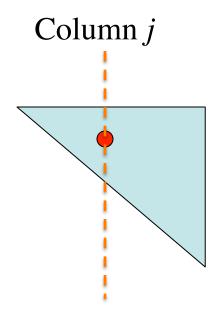
$$= A - (L_{t-1}P_{t-1}L_{t-2}P_{t-2}\cdots L_0P_0)^{-1}\lambda e_i e_j^T = A - de_j^T$$

$$\tilde{P}\left[\tilde{A}, A \times e, A \times w\right] = \tilde{L}\left[\tilde{U}, \tilde{c}, \tilde{v}\right], \quad \tilde{A} = A + de_{j}^{T}$$

$$\Rightarrow \tilde{P}\left[\tilde{A}, Ae, Aw\right] = \tilde{L}\left[\tilde{U}, \tilde{c}, \tilde{v}\right]$$

$$\Rightarrow \begin{cases} \tilde{P}\tilde{A} = \tilde{L}\tilde{U} \\ \tilde{P}Ae = \tilde{L}\tilde{c} \\ \tilde{P}Aw = \tilde{L}\tilde{v} \end{cases}$$

$$G = \begin{bmatrix} e^T \\ w^T \end{bmatrix} \begin{bmatrix} 1 & 1 & \cdots & 1 \\ w_1 & w_2 & \cdots & w_n \end{bmatrix}^T$$



- Luk's work
- Sherman Morison Formula

$$(A + uv^{T})^{-1} = A^{-1} - \frac{A^{-1}uv^{T}A^{-1}}{1 + v^{T}A^{-1}u}.$$

Given:

$$\begin{cases} \tilde{P}\tilde{A} = \tilde{L}\tilde{U} \\ \tilde{A}\tilde{x} = b \end{cases}$$

To Solve:

$$Ax = b$$

$$Ax = b$$

$$\Rightarrow x = A^{-1}b$$

$$\Rightarrow x = A^{-1}(\tilde{P}^{-1}\tilde{P})b = (\tilde{P}A)^{-1}\tilde{P}b$$

$$(\tilde{P}A)^{-1} = ?$$

Recall:

$$A - \tilde{A} = de_j^T$$

Therefore:

$$\begin{split} \tilde{P}A - \tilde{P}\tilde{A} &= (\tilde{P}a_{\cdot j} - \tilde{L}\tilde{U}_{\cdot j})e_{j}^{T} \\ \tilde{P}A &= \tilde{L}\tilde{U} + \tilde{L}(\tilde{L}^{-1}\tilde{P}a_{\cdot j} - \tilde{U}_{\cdot j})e_{j}^{T} = \tilde{L}(\tilde{U} + te_{j}^{T}) \\ &= \tilde{L}\tilde{U}(I + \tilde{U}^{-1}te_{j}^{T}) = \tilde{L}\tilde{U}(I + ve_{j}^{T}) \\ t &= \tilde{L}^{-1}\tilde{P}a_{\cdot j} - \tilde{U}_{\cdot j} \\ v &= \tilde{U}^{-1}t \end{split}$$

$$(\tilde{P}A)^{-1} = (\tilde{L}\tilde{U}(I + ve_j^T))$$

$$= (I + ve_j^T)^{-1}(\tilde{L}\tilde{U})^{-1}$$
Sherman
Morrison
$$= \left(I - \frac{1}{1 + v_j}ve_j^T\right)(\tilde{L}\tilde{U})^{-1}$$

$$Ax = b$$

$$= \left(I - \frac{1}{1 + v_j} v e_j^T\right) \tilde{x}$$

$$(1) \quad \tilde{L}\tilde{U}\tilde{x} = \tilde{P}b$$

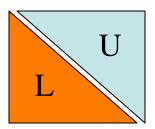
$$\begin{cases}
t = \tilde{L}^{-1}\tilde{P}a_{\cdot j} - \tilde{U}_{\cdot j} \\
v = \tilde{U}^{-1}t
\end{cases}$$

$$x = \left(I - \frac{y_j}{1 + v_j} v e_j^T\right) \tilde{x}$$

(1)
$$\tilde{L}\tilde{U}\tilde{x} = \tilde{P}b$$
 Needs protection
$$\begin{cases}
t = \tilde{L}^{-1}\tilde{P}a_{\cdot j} - \tilde{U}_{\cdot j} \\
v = \tilde{U}^{-1}t
\end{cases}$$
(2)
$$\begin{cases}
x = \left(I - \frac{y_j}{1 + v_j}ve_j^T\right)\tilde{x}
\end{cases}$$

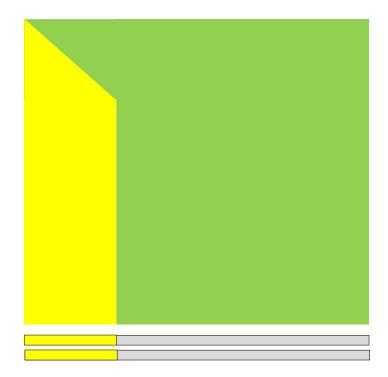
How to detect & recovery a soft error in L?

- The recovery of Ax=b requires a correct L
- L does not change once produced
 - Static checkpointing for L
- Delay pivoting on L to prevent checksum of L from being invalidated



Checkpointing for L, idea 1

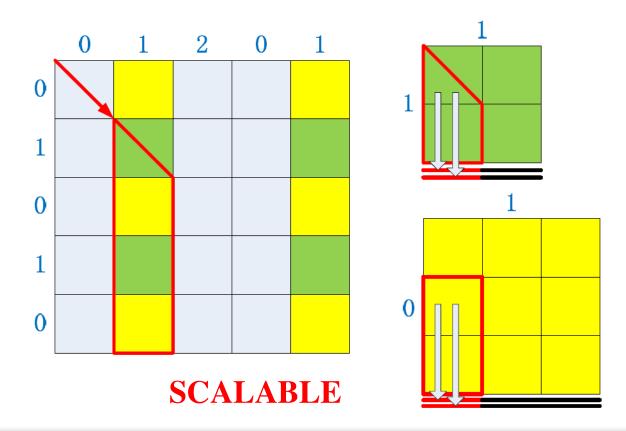
- PDGEMM based checkpointing
- Checkpointing time increases when scaled to more processes and larger matrices



NOT SCALABLE

Checkpointing for L, idea 2

- Local Checkpointing
- Each process checkpoints their local involved data
- Constant checkpointing time



Encoding for L

• On each process, for a column of L $l = [l_1, l_2, \dots, l_n]$

$$\begin{cases} l_1 + l_2 + \dots + l_n = c_1 \\ w_1 l_1 + w_2 l_2 + \dots + w_n l_n = c_2 \end{cases}$$

$$\begin{cases} l_1 + \dots + \tilde{l}_1 + \dots + l_n = \tilde{c}_1 \end{cases}$$

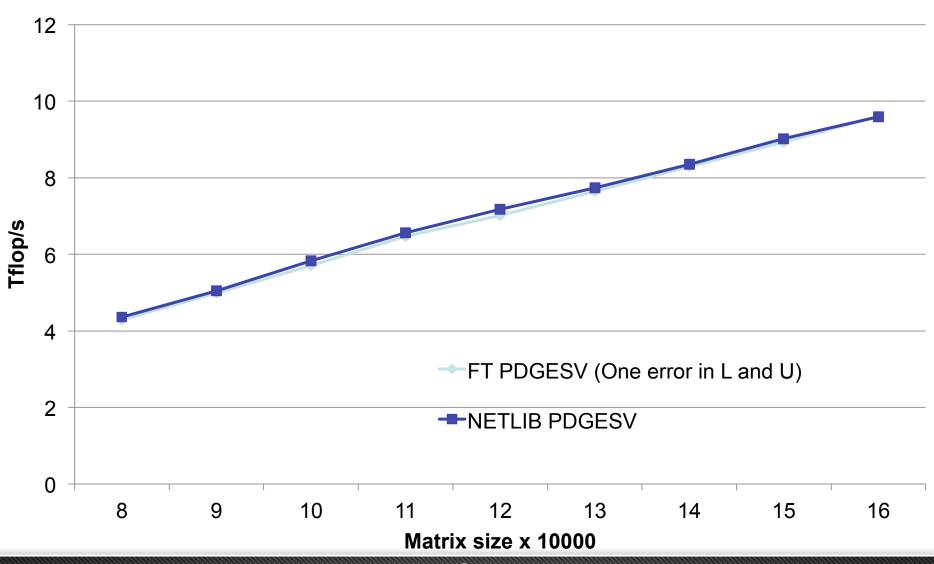
$$\begin{cases} l_1 + \dots + \tilde{l}_i + \dots + l_n = \tilde{c}_1 \\ w_1 l_1 + \dots + w_i \tilde{l}_i + \dots + w_n l_n = \tilde{c}_2 \end{cases}$$

$$\begin{cases} c_{1} - \tilde{c}_{1} = l_{i} - \tilde{l}_{i} \\ c_{2} - \tilde{c}_{2} = w_{i}(l_{i} - \tilde{l}_{i}) \end{cases} \qquad w_{i} = \frac{c_{2} - \tilde{c}_{2}}{c_{1} - \tilde{c}_{1}}$$

Kraken Performance

Two 2.6 GHz six-core AMD Opteron processors per node

32x32 MPI processes, **6** threads/(process, core) **6,144** cores used in total

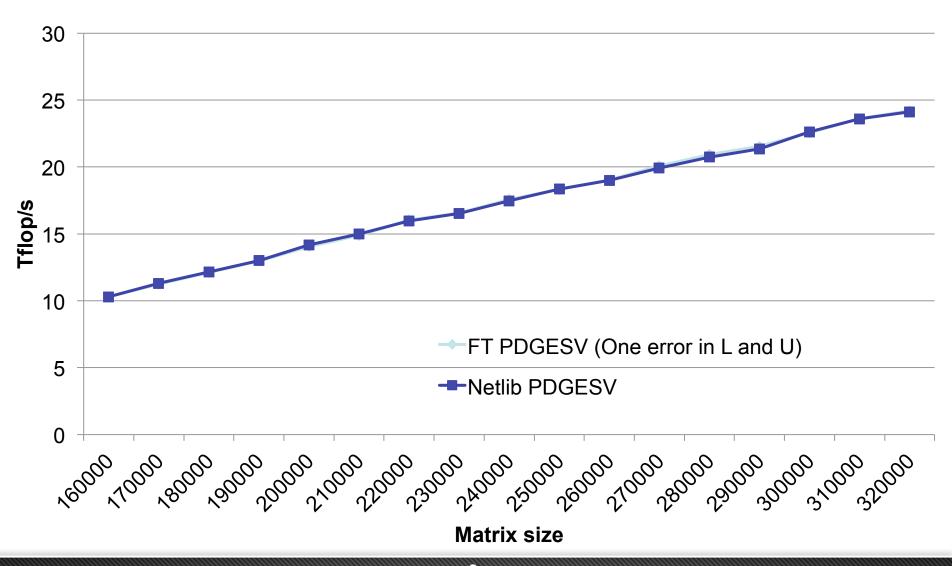




Kraken Performance

Two 2.6 GHz six-core AMD Opteron processors per node

64x64 MPI processes, **6** threads/(process, core) **24,576** used cores in total



Quastion?

Backup slides

September 29, 2011

$$\tilde{P}Ae = \tilde{L}\tilde{c}$$

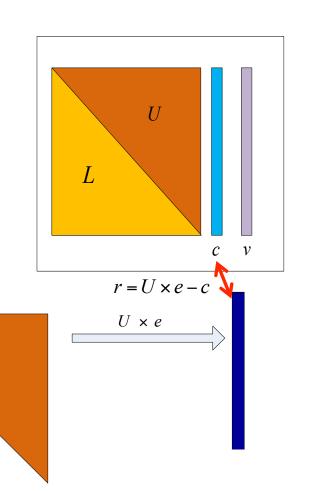
$$\Rightarrow \quad \tilde{c} = \tilde{L}^{-1}\tilde{P}Ae = \tilde{L}^{-1}\tilde{P}(\tilde{A} + de_{j}^{T})e$$

$$= \tilde{L}^{-1}(\tilde{P}\tilde{A} + \tilde{P}de_{j}^{T})e$$

$$= \tilde{L}^{-1}(\tilde{L}\tilde{U} + \tilde{P}de_{j}^{T})e$$

$$= \tilde{U}e + \tilde{L}^{-1}\tilde{P}d$$

$$\Rightarrow \quad \tilde{c} - \tilde{U}e = \tilde{L}^{-1}\tilde{P}d = r$$



$$\tilde{P}Aw = \tilde{L}\tilde{v}$$

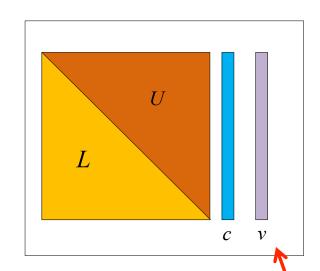
$$\Rightarrow \quad \tilde{\mathbf{v}} = \tilde{L}^{-1}\tilde{P}Aw = \tilde{L}^{-1}\tilde{P}(\tilde{A} + de_{j}^{T})w$$

$$= \tilde{L}^{-1}(\tilde{P}\tilde{A} + \tilde{P}de_{j}^{T})w$$

$$= \tilde{L}^{-1}(\tilde{L}\tilde{U} + \tilde{P}de_{j}^{T})w$$

$$= \tilde{U}w + \tilde{L}^{-1}\tilde{P}dw_{j}$$

$$\Rightarrow \quad \tilde{\mathbf{v}} - \tilde{U}w = \tilde{L}^{-1}\tilde{P}dw_{j} = s$$



 $s = U \times w - v$

$$\begin{cases} \tilde{c} - \tilde{U}e = \tilde{L}^{-1}\tilde{P}d = r \\ \tilde{v} - \tilde{U}w = w_{j}\tilde{L}^{-1}\tilde{P}d = s \end{cases} \Rightarrow S = w_{j} \times r$$

$$\Rightarrow \mathbf{w}_{j} \begin{bmatrix} 1 \\ 1 \\ \vdots \\ 1 \end{bmatrix} = s./r$$

- W_j is the j_{th} element of vector w in the generator matrix
- Component-wise division of s and r reveals w_j
- Search w_i in w reveals the **initial soft error's column**

Extra Storage

- For input matrix of size MxN on PxQ grid
 - A copy of the original matrix
 - Not necessary when it's easy to re-generate the required column of the original matrix
 - 2 additional columns: 2 x M
 Each process has 2 rows: $2 \times \frac{N}{Q}$, in total $P \times 2 \times N$

$$r = \frac{extra\ storage}{matrix\ storage} = \frac{2 \times M + P \times 2 \times N}{M \times N}$$
$$= \frac{2}{N} + \frac{P \times 2}{M} \xrightarrow{N \to \infty} \frac{P \times 2}{M}$$