# High Performance Dense Linear System Solver with Soft Error Resilience 

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## Agenda

- Soft error threat to the dense linear solver
- LU factorization
- Error propagation
- Error modeling
- Fault tolerant algorithm
- Performance Evaluation


## Soft error

- Silent error due to radiation
- Alpha particle
- High energy neutron
- Thermal neutron

- Outbreaks
- Commercial computing system from Sun Microsystem in 2000
- ASC Q supercomputer at Los Alamos National Lab in 2003


## LU based linear solver

$$
\begin{aligned}
& A x=b \\
& A=L U \\
& x=U \backslash(L \backslash b)
\end{aligned}
$$

## Block LU factorization



GEMM

GETF2


GETF2


TRSM


STRSM

## General work flow

(1) Generate checksum for the input matrix as additional columns
(2) Perform LU factorization WITH the additional checksum columns
(3) Solve $A x=b$ using LU from the factorization (even if soft error occurs during LU factorization)
(4) Check for soft error
(5) Correct solution $x$


## Why is soft error hard to handle?

- Soft error occurs silently
- Propagation


## Example: Error propagation

## Error location (using matlab notation and 1based index)

Error strikes right before panel factorization of (41:200, 41:60),

Case 1: Error at $(35,10)$, in $L$ area

Case 2: Error at $(50,120)$, in A' area

Note: Pivoting on the left of panel factorization is delayed to the end of error detection and recovery so that error in $L$ area does not move

## Case 1: Non-propagating error



## Case 2: Propagating error



## Soft error challenge

## When?

## Where?

## Error modeling (for propagating error)

- When?
- Answer: Doesn't really matter



## Error modeling (for "where")

Input matrix $A$
One step of LU $\quad A_{t}=L_{t-1} P_{t-1} A_{t-1}$
If no soft error occurs $U=\left(L_{n} P_{n}\right) \cdots\left(L_{1} P_{1}\right)\left(L_{0} P\right)_{0} A_{0}$
If soft error occurs at step t $\tilde{A}_{t}=L_{t-1} P_{t-1} A_{t-1}-\lambda e_{i} e_{j}^{T}$

$$
=L_{t-1} P_{t-1}\left(L_{t-2} P_{t-2} \cdots L_{0} P_{0}\right) A_{0}-\lambda e_{i} e_{j}^{T}
$$

Define an initial erroneous initial matrix $\tilde{A}$

$$
\begin{aligned}
& \tilde{A} \cong\left(L_{t-1} P_{t-1} L_{t-2} P_{t-2} \cdots L_{0} P_{0}\right)^{-1} \tilde{A}_{t} \\
& =A-\left(L_{t-1} P_{t-1} L_{t-2} P_{t-2} \cdots L_{0} P_{0}\right)^{-1} \lambda e_{i} e_{j}^{T}=A-d e_{j}^{T}
\end{aligned}
$$

## Locate Error

$$
\begin{aligned}
& \tilde{P}[\tilde{A}, A \times e, A \times w]=\tilde{L}[\tilde{U}, \tilde{c}, \tilde{v}], \quad \tilde{A}=A+d e_{j}^{T} \\
& \Rightarrow \tilde{P}[\tilde{A}, A e, A w]=\tilde{L}[\tilde{U}, \tilde{c}, \tilde{v}] \\
& \Rightarrow\left\{\begin{array}{l}
\tilde{P} \tilde{A}=\tilde{L} \tilde{U} \\
\tilde{P} A e=\tilde{L} \tilde{C} \\
\tilde{P} A w=\tilde{L} \tilde{v}
\end{array}\right. \\
& G=\left[\begin{array}{c}
e^{T} \\
w^{T}
\end{array}\right]\left[\begin{array}{cccc}
1 & 1 & \cdots & 1 \\
w_{1} & w_{2} & \cdots & w_{n}
\end{array}\right]^{T} \\
& \text { Column } j
\end{aligned}
$$

## Recover Ax=b

- Luk's work
- Sherman Morison Formula

$$
\left(A+u v^{T}\right)^{-1}=A^{-1}-\frac{A^{-1} u v^{T} A^{-1}}{1+v^{T} A^{-1} u}
$$

## Recover Ax=b

Given:

$$
\left\{\begin{array}{c}
\tilde{P} \tilde{A}=\tilde{L} \tilde{U} \\
\tilde{A} \tilde{x}=b
\end{array}\right.
$$

To Solve:

$$
A x=b
$$

## Recover Ax=b

$$
\begin{aligned}
& A x=b \\
& \Rightarrow \quad x=A^{-1} b \\
& \Rightarrow \quad x=A^{-1}\left(\tilde{P}^{-1} \tilde{P}\right) b=(\tilde{P} A)^{-1} \tilde{P} b \\
& (\tilde{P} A)^{-1}=?
\end{aligned}
$$

## Recover Ax=b

Recall:

$$
A-\tilde{A}=d e_{j}^{T}
$$

Therefore:

$$
\begin{aligned}
& \quad \tilde{P} A-\tilde{P} \tilde{A}=\left(\tilde{P} a_{\cdot j}-\tilde{L} \tilde{U}_{\cdot j}\right) e_{j}^{T} \\
& \tilde{P} A=\tilde{L} \tilde{U}+\tilde{L}\left(\tilde{L}^{-1} \tilde{P} a_{\cdot j}-\tilde{U}_{\cdot j}\right) e_{j}^{T}=\tilde{L}\left(\tilde{U}+t e_{j}^{T}\right) \\
& =\tilde{L} \tilde{U}\left(I+\tilde{U}^{-1} t e_{j}^{T}\right)=\tilde{L} \tilde{U}\left(I+v e_{j}^{T}\right) \\
& \boldsymbol{t}=\tilde{L}^{-1} \tilde{P} a_{\cdot j}-\tilde{U}_{\cdot j} \\
& \boldsymbol{v}=\tilde{U}^{-1} t
\end{aligned}
$$

## Recover Ax=b

Sherman
Morrison

$$
\begin{aligned}
& (\tilde{P} A)^{-1}=\left(\tilde{L} \tilde{U}\left(I+v e_{j}^{T}\right)\right) \\
& =\left(I+v e_{j}^{T}\right)^{-1}(\tilde{L} \tilde{U})^{-1} \\
& =\left(I-\frac{1}{1+v_{j}} v e_{j}^{T}\right)(\tilde{L} \tilde{U})^{-1}
\end{aligned}
$$

## Recover Ax=b

$$
\begin{aligned}
& A x=b \\
& =\left(I-\frac{1}{1+v_{j}} v e_{j}^{T}\right) \tilde{x}
\end{aligned}
$$

## Recover Ax=b

(1) $\tilde{L} \tilde{U} \tilde{x}=\tilde{P} b$
(2) $\left\{\begin{array}{c}t=\tilde{L}^{-1} \tilde{P} a_{\cdot j}-\tilde{U}_{\cdot j} \\ v=\tilde{U}^{-1} t \\ x=\left(I-\frac{y_{j}}{1+v_{j}} v e_{j}^{T}\right) \tilde{x}\end{array}\right.$

## Recover Ax=b

(1) $\tilde{L} \tilde{U} \tilde{x}=\tilde{P} b$
$t=\tilde{L}^{-1} \tilde{P} a_{\cdot j}-\tilde{U}_{\cdot j}$
$v=\tilde{U}^{-1} t$
$x=\left(I-\frac{y_{j}}{1+v_{j}} v e_{j}^{T}\right) \tilde{x}$

## How to detect \& recovery a soft error in L?

- The recovery of $A x=b$ requires a correct $L$
- L does not change once produced
- Static checkpointing for L
- Delay pivoting on $L$ to prevent checksum of $L$ from being invalidated



## Checkpointing for L, idea 1

- PDGEMM based checkpointing
- Checkpointing time increases when scaled to more processes and larger matrices

NOT SCALABLE

Ichion

## Checkpointing for L, idea 2

- Local Checkpointing
- Each process checkpoints their local involved data
- Constant checkpointing time



## Encoding for L

- On each process, for a column of $\mathrm{L} \quad l=\left[l_{1}, l_{2}, \cdots, l_{n}\right]$

$$
\begin{gathered}
\left\{\begin{array}{c}
l_{1}+l_{2}+\cdots+l_{n}=c_{1} \\
w_{1} l_{1}+w_{2} l_{2}+\cdots+w_{n} l_{n}=c_{2}
\end{array}\right. \\
\left\{\begin{array}{c}
l_{1}+\cdots+\tilde{l}_{i}+\cdots+l_{n}=\tilde{c}_{1} \\
w_{1} l_{1}+\cdots+w_{i} \tilde{l}_{i}+\cdots+w_{n} l_{n}=\tilde{c}_{2}
\end{array}\right. \\
\left\{\begin{array}{c}
c_{1}-\tilde{c}_{1}=l_{i}-\tilde{l}_{i} \\
c_{2}-\tilde{c}_{2}=w_{i}\left(l_{i}-\tilde{l}_{i}\right)
\end{array} w_{i}=\frac{c_{2}-\tilde{c}_{2}}{c_{1}-\tilde{c}_{1}}\right.
\end{gathered}
$$

## Kraken Performance

Two 2.6 GHz six-core AMD Opteron processors per node
32x32 MPI processes, 6 threads/(process, core)
6,144 cores used in total


## Kraken

Two 2.6 GHz six-core AMD Opteron processors per node
64x64 MPI processes, 6 threads/(process, core) 24,576 used cores in total


## © (2xనtiom?

- Backup slides


## Locate Error

$$
\begin{aligned}
& \tilde{P} A e=\tilde{L} \tilde{c} \\
& \begin{aligned}
& \Rightarrow \quad \tilde{c}= \tilde{L}^{-1} \tilde{P} A e=\tilde{L}^{-1} \tilde{P}\left(\tilde{A}+d e_{j}^{T}\right) e \\
&=\tilde{L}^{-1}\left(\tilde{P} \tilde{A}+\tilde{P} d e_{j}^{T}\right) e \\
&=\tilde{L}^{-1}\left(\tilde{L} \tilde{U}+\tilde{P} d e_{j}^{T}\right) e \\
&=\tilde{U} e+\tilde{L}^{-1} \tilde{P} d \\
& \Rightarrow \quad \tilde{c}-\tilde{U} e=\tilde{L}^{-1} \tilde{P} d=r
\end{aligned}
\end{aligned}
$$



## Locate Error

$$
\begin{aligned}
& \tilde{P} A w= \\
& \begin{aligned}
\Rightarrow \quad \tilde{L} \tilde{v} & =\tilde{L}^{-1} \tilde{P} A w=\tilde{L}^{-1} \tilde{P}\left(\tilde{A}+d e_{j}^{T}\right) w \\
& =\tilde{L}^{-1}\left(\tilde{P} \tilde{A}+\tilde{P} d e_{j}^{T}\right) w \\
& =\tilde{L}^{-1}\left(\tilde{L} \tilde{U}+\tilde{P} d e_{j}^{T}\right) w \\
& =\tilde{U} w+\tilde{L}^{-1} \tilde{P} d w_{j} \\
\Rightarrow \quad \tilde{v} & -\tilde{U} w=\tilde{L}^{-1} \tilde{P} d w_{j}=s
\end{aligned}
\end{aligned}
$$



## Locate Error

$$
\begin{aligned}
& \left\{\begin{array}{c}
\tilde{c}-\tilde{U} e=\tilde{L}^{-1} \tilde{P} d=r \\
\tilde{v}-\tilde{U} w=w_{j} \tilde{L}^{-1} \tilde{P} d=s
\end{array} \Rightarrow s=w_{j} \times r\right. \\
& \Rightarrow w_{j}\left[\begin{array}{c}
1 \\
1 \\
\vdots \\
1
\end{array}\right]=s . / r
\end{aligned}
$$

- $W_{j}$ is the $\mathrm{j}_{\mathrm{th}}$ element of vector w in the generator matrix
- Component-wise division of $s$ and $r$ reveals $w_{j}$
- Search $w_{j}$ in w reveals the initial soft error's column


## Extra Storage

- For input matrix of size $M x N$ on $P x Q$ grid
- A copy of the original matrix
- Not necessary when it's easy to re-generate the required column of the original matrix
- 2 additional columns: $2 \times \mathrm{M}$
- Each process has 2 rows: $2 \times \frac{N}{Q}$, in total $P \times 2 \times N$

$$
\begin{aligned}
& r=\frac{\text { extra storage }}{\text { matrix storage }}=\frac{2 \times M+P \times 2 \times N}{M \times N} \\
& =\frac{2}{N}+\frac{P \times 2}{M} \xrightarrow{N \rightarrow \infty} \frac{P \times 2}{M}
\end{aligned}
$$

