

Polarities, quasi-symmetric designs, and Hamada's conjecture

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1 **Abstract** We prove that every polarity of $PG(2k - 1, q)$, where $k \geq 2$, gives rise to a
2 design with the same parameters and the same intersection numbers as, but not isomorphic to,
3 $PG_k(2k, q)$. In particular, the case $k = 2$ yields a new family of quasi-symmetric designs. We
4 also show that our construction provides an infinite family of counterexamples to Hamada's
5 conjecture, for any field of prime order p . Previously, only a handful of counterexamples
6 were known.

7 **Keywords** Polarity · Projective geometry · Design · Quasi-symmetric design · Hamada's
8 conjecture

9 **Mathematics Subject Classification (2000)** ■

10 1 Introduction

11 We prove that every polarity of $PG(2k - 1, q)$, where $k \geq 2$, gives rise to a design with the
12 same parameters and the same intersection numbers as, but not isomorphic to, $PG_k(2k, q)$,
13 the design of points and k -spaces in projective $2k$ -space over $GF(q)$. The new designs are
14 obtained by distorting the classical geometric designs with the help of the given polarity,
15 acting on a fixed hyperplane in $PG(2k, q)$. In particular, the case $k = 2$ yields a new family
16 of quasi-symmetric designs.

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17 By construction, our new examples of designs with classical geometric parameters still
 18 share many properties with the geometric designs $PG_k(2k, q)$. In particular, there always is
 19 a set H of $q^{2k-1} + \dots + q + 1$ points on which the blocks of the design induce an isomorphic
 20 copy of $PG(2k - 1, q)$, while a copy of an affine $2k$ -space $AG(2k, q)$ is induced on the set
 21 A formed by the remaining q^{2k} points. Moreover, the lines of the design joining two points
 22 of H or two points of A still have the natural geometric size, that is, $q + 1$ or q , respectively,
 23 whereas a point of H and a point of A always determine a line of size 2.

24 We also show that our construction provides an infinite family of counterexamples to
 25 Hamada's conjecture [8] from 1973, for fields of arbitrary prime order, and for any dimension
 26 $2k \geq 4$. Previously, only a handful of counterexamples were known, namely two parameter
 27 sets, 2-(31, 7, 7) [7,29] and 3-(32, 8, 7) [29] for the binary case, and a single parameter set
 28 2-(64, 16, 5) for the quaternary case ($q = 4$) [10,22].

29 Hamada's conjecture is of fundamental importance for two reasons. First, it indicates
 30 that the classical geometric designs, as designs having minimum p -rank among all possi-
 31 ble designs with the given parameters, are the best choice to use for the construction
 32 of error-correcting codes with majority-logic decoding [24,25]. It is known that the
 33 number of non-isomorphic designs having the same parameters as the classical geomet-
 34 ric designs of hyperplanes in $PG(n, q)$ or $AG(n, q)$, $n \geq 3$, grows exponentially with linear
 35 growth of n [15,17–19]. Secondly, the conjecture provides an elegant and computationally
 36 simple characterization of the classical geometric designs in terms of the p -rank of their inci-
 37 dence matrices: the complexity of computing the rank of a matrix is a cubic polynomial in the
 38 number of rows (or columns), while the complexity of finding isomorphisms between block
 39 designs is as hard as the notoriously difficult graph isomorphism problem; see [5, Remark
 40 VII.6.6].

41 2 A construction method for pseudo-geometric designs

42 We begin by describing a general method for constructing 2-designs with the same parameters
 43 as some classical geometric designs. To this end, let Π denote $PG(2k, q)$, the $2k$ -dimen-
 44 sional projective space over the field $GF(q)$ with q elements. As is well-known, the points
 45 and k -spaces of Π form a 2- (v, K, λ) design $\mathcal{D} = PG_k(2k, q)$ with parameters

$$46 \quad v = \frac{q^{2k+1} - 1}{q - 1}, \quad K = \frac{q^{k+1} - 1}{q - 1}, \quad \lambda = \frac{(q^{2k-1} - 1) \dots (q^{k+1} - 1)}{(q^{k-1} - 1) \dots (q - 1)}. \quad (1)$$

47 Furthermore, the lines of \mathcal{D} are just the lines of Π and hence all have cardinality $q + 1$.¹

48 Now let H denote a fixed hyperplane of Π . Trivially, the subspaces of Π induce a geom-
 49 etry Π_0 isomorphic to $PG(2k - 1, q)$ on H . Since the lines of Π_0 are just those lines of \mathcal{D}
 50 which are contained in H , we may view H as a copy of the projective space $PG(2k - 1, q)$
 51 in the design \mathcal{D} . Moreover, let A be the set of points not contained in H . Then the subspaces
 52 of Π induce a geometry Σ isomorphic to the affine space $AG(2k, q)$ on A . Now each line
 53 of Σ corresponds to a line ℓ of Π ; of course, considered as a line of the affine space Σ , the
 54 projective line ℓ loses its *infinite point* $\ell \cap H$. Since ℓ is also a line of the design \mathcal{D} , we may
 55 view A as a copy of the affine space $AG(2k, q)$ in \mathcal{D} . In view of the preceding observations,
 56 we shall refer to H also as a *hyperplane* of \mathcal{D} .

¹ Recall that the *line* determined by two points of a design is defined as the intersection of all blocks containing these two points. See [4] for background on designs, and [12,11] for background on finite projective spaces.

57 More generally, let \mathcal{D}' be any design with the parameters (1) of $PG_k(2k, q)$. If there exists
 58 a set H of $q^{2k-1} + \dots + q + 1$ points on which the lines of the design induce an isomorphic
 59 copy of $PG(2k - 1, q)$, while a copy of $AG(2k, q)$ is induced on the set A formed by the
 60 remaining q^{2k} points, we shall call H a *hyperplane* of \mathcal{D}' . The points of H will be referred
 61 to as the *infinite points* of \mathcal{D}' , and the points in A as the *affine points* of \mathcal{D}' .

62 If, in addition, the intersection numbers of \mathcal{D}' are the same as those of $PG_k(2k, q)$, we
 63 shall call \mathcal{D}' a *pseudo-geometric* design. Our main result will be a construction method for
 64 designs which are pseudo-geometric but not actually geometric.

65 We begin with a more general method yielding designs with the parameters of a geometric
 66 design $PG_k(2k, q)$. Given any *affine* block B of $\mathcal{D} = PG_k(2k, q)$ —that is, any k -space of
 67 Π which is not contained in the hyperplane H —we write B in the form

$$68 \quad B = B_\infty \cup B_{\text{aff}}, \quad (2)$$

69 where $B_\infty := B \cap H$ is a projective $(k - 1)$ -space contained in the hyperplane H and
 70 $B_{\text{aff}} := B \cap A$ is a k -space of the affine space Σ induced on A . In particular,

$$71 \quad |B_{\text{aff}}| = q^k \quad \text{and} \quad |B_\infty| = q^{k-1} + \dots + q + 1 = \frac{q^k - 1}{q - 1}. \quad (3)$$

72 If B and C are affine blocks and $B_\infty = C_\infty$, then B_{aff} and C_{aff} are affine translates of each
 73 other.

74 Now let α be any permutation of the projective $(k - 1)$ -spaces contained in H , and associate
 75 with each affine block B of \mathcal{D} a point set $\alpha(B)$ as follows:

$$76 \quad \alpha(B) := \alpha(B_\infty) \cup B_{\text{aff}}. \quad (4)$$

77 Thus, we keep the affine points of all affine blocks unchanged, and merely exchange their
 78 infinite parts, using the permutation α . We shall denote the incidence structure obtained from
 79 \mathcal{D} by replacing each affine block B by its distorted version $\alpha(B)$ as $\alpha(\mathcal{D})$. Then it is easy to
 80 prove the following result.

81 **Lemma 2.1** *For each permutation α of the $(k - 1)$ -spaces contained in H , the incidence*
 82 *structure $\alpha(\mathcal{D})$ is a 2-design with the same parameters as $\mathcal{D} = PG_k(2k, q)$.*

83 In general, the designs just constructed may have intersection numbers different from
 84 those of \mathcal{D} . If we wish to preserve intersection sizes, we will have to choose α judiciously.
 85 Before we address this problem, let us remark that Lemma 2.1 can be used to show that the
 86 number of 2-designs with the parameters of $PG_k(2k, q)$ grows exponentially; this is a special
 87 case of a more general result which will be presented elsewhere [16].

88 As it turns out, our aim can be achieved by choosing α as a polarity of the projective space
 89 $\Pi_0 \cong PG(2k - 1, q)$ induced on H . Recall that a *polarity* of a projective space $PG(n, q)$ is
 90 an involutory isomorphism between $PG(n, q)$ and its dual space; in other words, a polarity
 91 is an incidence preserving bijection interchanging points and hyperplanes. Note that any
 92 polarity of Π_0 maps i -spaces to $(2k - i - 2)$ -spaces, for $i = 0, \dots, 2k - 2$; in particular,
 93 α induces a permutation on the $(k - 1)$ -spaces contained in H , and hence can be used in
 94 our construction. We refer the reader to [12] for a thorough discussion of polarities in finite
 95 projective spaces.

96 **Lemma 2.2** *For each polarity α of $\Pi_0 \cong PG(2k - 1, q)$, the design $\alpha(\mathcal{D})$ has the same*
 97 *intersection numbers as $\mathcal{D} = PG_k(2k, q)$.*

98 *Proof* The interesting case to consider concerns the intersection sizes of blocks of \mathcal{D}' which
 99 correspond to affine blocks of \mathcal{D} . As we will see, α even preserves these intersection sizes:

$$100 \quad |\alpha(B) \cap \alpha(C)| = |B \cap C| \quad (5)$$

101 for any two affine blocks B and C of \mathcal{D} .

102 With the notation introduced in (2), the infinite parts B_∞ and C_∞ of the given two blocks
 103 are $(k - 1)$ -subspaces of $H \cong PG(2k - 1, q)$. In view of the construction given in (4), the
 104 validity of (5) is clear provided that $B_\infty = C_\infty$.

105 Next, note that B_∞ and C_∞ are disjoint if and only if their images under α are disjoint.
 106 Indeed, by the dimension formula, these two $(k - 1)$ -subspaces of H intersect if and only
 107 if they are both contained in a hyperplane H_0 of H ; as α is incidence preserving, this holds
 108 if and only if their images $\alpha(B_\infty)$ and $\alpha(C_\infty)$ intersect in the point $\alpha(H_0)$. This proves the
 109 validity of (5) in the special case where $B_\infty \cap C_\infty = \emptyset$.

110 We may now assume that $U := B_\infty \cap C_\infty$ is an i -subspace of H , where $0 \leq i \leq k - 2$. Then
 111 $\alpha(U)$ is a $(2k - i - 2)$ -subspace which contains the two $(k - 1)$ -spaces $\alpha(B_\infty)$ and $\alpha(C_\infty)$,
 112 as α is incidence preserving. Again using the dimension formula, $\alpha(B_\infty)$ and $\alpha(C_\infty)$ have
 113 to intersect in a j -subspace for some $j \geq i$. Applying this argument to $\alpha(B_\infty)$ and $\alpha(C_\infty)$
 114 and using that α is an involution shows that also $i \geq j$. Hence $\alpha(B_\infty) \cap \alpha(C_\infty)$ is again an
 115 i -subspace, and therefore (5) holds also in the case $B_\infty \cap C_\infty \neq \emptyset$.

116 Finally, note that the multiset of the remaining intersection numbers does not change, as
 117 blocks of \mathcal{D} contained in H are kept in $\alpha(\mathcal{D})$ and as the infinite parts of the affine blocks are
 118 merely permuted under α . (However, in general, the image $\alpha(B)$ of a given affine block B
 119 may intersect a specific infinite block C in a different manner as B does). \square

120 **Lemma 2.3** For each polarity α of $\Pi_0 \cong PG(2k - 1, q)$, the design $\alpha(\mathcal{D})$ has line sizes
 121 $q + 1$, q and 2. More precisely, any line of $\alpha(\mathcal{D})$ joining two infinite points has cardinality
 122 $q + 1$; any line of $\alpha(\mathcal{D})$ joining two affine points has cardinality q ; finally, an infinite point
 123 and an affine point always determine a line of size 2 in $\alpha(\mathcal{D})$.

124 *Proof* Let us consider a fixed (affine) $(k - 1)$ -subspace U_{aff} of the affine space $\Sigma \cong$
 125 $AG(2k, q)$ induced on the set A of affine points of \mathcal{D} . Then U_{aff} is contained in exactly
 126 $q^k + \dots + q + 1$ affine blocks of \mathcal{D} , as this is the number of k -dimensional subspaces of
 127 $AG(2k, q)$ containing a given $(k - 1)$ -space. Recall that each such block B has the form
 128 given in (2).

129 Now U_{aff} extends to a unique $(k - 1)$ -subspace U of the underlying projective space Π .
 130 Note that U contains exactly $q^{k-2} + \dots + q + 1$ infinite points, as U has to intersect the
 131 hyperplane H of Π in a $(k - 2)$ -dimensional subspace U_∞ . Hence any two distinct affine
 132 blocks containing U_{aff} share exactly $q^{k-2} + \dots + q + 1$ infinite points, namely those in U_∞ ;
 133 and by (3), any such block B has precisely q^{k-1} infinite points outside of U_∞ . But then the
 134 remaining $q^{2k-1} + \dots + q^k + q^{k-1}$ infinite points are partitioned by the $q^k + \dots + q + 1$ affine
 135 blocks through U_{aff} :

$$136 \quad (q^k + \dots + q + 1)q^{k-1} = q^{2k-1} + \dots + q^k + q^{k-1}.$$

137 Thus, the infinite parts B_∞ of the $q^k + \dots + q + 1$ affine blocks B through U_{aff} form the bundle
 138 of $(k - 1)$ -subspaces of H through the common $(k - 2)$ -subspace U_∞ . Under the polarity
 139 α , this bundle is mapped to a set of $q^k + \dots + q + 1$ $(k - 1)$ -dimensional subspaces of the
 140 k -subspace $\alpha(U_\infty)$. Hence, the images $\alpha(B_\infty)$ are simply all hyperplanes of the projective
 141 space $\alpha(U_\infty) \cong PG(k, q)$. Therefore, the images of the infinite parts of any two distinct
 142 affine blocks through U_{aff} intersect in a $(k - 2)$ -dimensional subspace of $\alpha(U_\infty)$. Hence,

no point of U_∞ lies in the intersection of all affine blocks $\alpha(B)$ through U_{aff} , and thus the intersection of all these blocks in $\alpha(\mathcal{D})$ is simply U_{aff} .

Now, let ℓ be any line of \mathcal{D} joining two affine points, so that ℓ has size $q + 1$ and consists of q affine points and one infinite point ℓ_∞ . Note that ℓ is the intersection of all $(k - 1)$ -dimensional affine subspaces U_{aff} of Σ extended to subspaces U of Π , and we have just seen that the affine part U_{aff} of each such subspace U is simply the intersection of all blocks of $\alpha(\mathcal{D})$ containing U_{aff} . This shows that the line corresponding to ℓ in $\alpha(\mathcal{D})$ is precisely the q -set $\ell \setminus \ell_\infty$: the distortion by α results in ℓ losing its infinite point.²

Finally, it is clear that any line joining two infinite points of \mathcal{D} remains a line of $\alpha(\mathcal{D})$. Now it easily follows that an infinite point and an affine point always determine a line of size 2 in $\alpha(\mathcal{D})$. \square

Combining the preceding three Lemmas, we obtain our main result:

Theorem 2.4 Consider the design $\mathcal{D} = PG_k(2k, q)$. Let H be a hyperplane of \mathcal{D} , and let A be the set of points not in H . In addition, let α be any polarity of the hyperplane $H \cong PG(2k - 1, q)$. Then the design $\alpha(\mathcal{D})$ defined above is a pseudo-geometric design with the same parameters as, but not isomorphic to, $PG_k(2k, q)$. \square

We conclude this section by pointing out that any two polarities of $\Pi_0 \cong PG(2k - 1, q)$ lead to isomorphic pseudo-geometric designs, even if the polarities are of different types. While this might seem surprising, it is in fact easy to prove: the product of two polarities is a collineation, hence any two polarities differ by a collineation only. Now it is easy to check that applying a non-trivial collineation β in our construction yields a design $\beta(\mathcal{D})$ different from, but isomorphic to, \mathcal{D} .

3 New quasi-symmetric designs

In this section, we consider the special case $k = 2$. Here the points and planes of $\Pi = PG(4, q)$ yield a 2-design which is *quasi-symmetric*; that is, it has just two intersection numbers, namely 1 and $q + 1$. Also, the lines of this design are just the lines of Π and hence all have cardinality $q + 1$.³

The designs $PG_2(4, q)$ form a well-known family of quasi-symmetric designs. They have been studied quite intensively, and several characterizations are available. To mention the most natural result, a quasi-symmetric design with the parameters of $PG_2(4, q)$ and intersection numbers 1 and $q + 1$ is classical if and only if all lines have size $q + 1$. This is due to Sane and Shrikhande [26], who also gave various other characterizations.

Theorem 2.4 specializes to the following construction for a new family of quasi-symmetric designs with the parameters of $PG_2(4, q)$:

Theorem 3.1 Consider the design $\mathcal{D} = PG_2(4, q)$, let H be a hyperplane of \mathcal{D} , and let A be the set of points not in H . In addition, let α be any polarity of the hyperplane $H \cong PG(3, q)$. Then the design $\alpha(\mathcal{D})$ defined in Sect. 2 is a pseudo-geometric quasi-symmetric design with the same parameters as, but not isomorphic to, $PG_2(4, q)$. \square

With the exception of the smallest case, i.e. $q = 2$, none of the designs in Theorem 3.1 was known previously; thus we indeed have a new infinite family of quasi-symmetric designs.

² More generally, all subspaces of Π of dimension at most $k - 1$ which are not contained in H can be recovered as suitable intersections of blocks of \mathcal{D} ; under α , the intersection of the corresponding distorted blocks no longer contains an infinite point and simply is the original affine part of the subspace.

³ See [27] for a monograph on quasi-symmetric designs.

183 By a result of Tonchev [29], there are exactly five quasi-symmetric 2—(31,7,7)-designs
 184 with intersection numbers 1 and 3; among these designs is, of course, the classical example
 185 $PG_2(4, 2)$. It is interesting to note that just one of the further four examples contains a hyper-
 186 plane; hence this design has to arise from Theorem 3.1. Actually, we discovered our general
 187 construction for pseudo-geometric designs when we tried to get a better understanding of
 188 this specific design which shares so many properties with the classical example. It seemed
 189 to us that there ought to be a geometric way of obtaining it—an intuition which fortunately
 190 turned out to be correct.

191 A more recent characterization of the geometric designs $PG_2(4, q)$ in terms of *good*
 192 *blocks*—a notion introduced in [23]—is due to Mavron, McDonough and Shrikhande [21].
 193 In any quasi-symmetric design with intersection numbers x and y , where $0 \leq x < y$, a
 194 block B is said to be *good* if, for any block C with $|B \cap C| = y$ and any point $p \notin C$, there
 195 is a (unique) block containing p and $B \cap C$. The result of [21] characterizes the geometric
 196 design $PG_2(4, q)$ among all quasi-symmetric designs with the same parameters and with
 197 intersection numbers 1 and $q + 1$ by the property that all blocks of the design are good.
 198 Subsequently, this result was strengthened by Baartmans and Sane [3] who proved that it
 199 suffices to assume that all the blocks passing through a fixed point p are good.

200 The authors of [21] also knew⁴ just one example of a quasi-symmetric design with the
 201 parameters of $PG_2(4, q)$ where some of the blocks, but not all blocks, are good, namely the
 202 pseudo-geometric 2—(31,7,7)-design discussed above. It is easy to check that if $\alpha(\mathcal{D})$ is a
 203 design obtained using a polarity α in a hyperplane H , then precisely the blocks contained in
 204 H are good.

205 4 Counterexamples to Hamada's conjecture

206 In this section, we shall see that our construction from Sect. 2 provides an infinite family of
 207 counterexamples to a famous conjecture by Hamada [8] from 1973. This conjecture reads as
 208 follows:

209 **Conjecture 4.1** (Hamada's Conjecture) *Let \mathcal{D} be a design with the parameters of a geomet-*
 210 *ric design $PG_d(n, q)$ or $AG_d(n, q)$, where q is a power of a prime p . Then the p -rank of the*
 211 *incidence matrix of \mathcal{D} is greater than or equal to the p -rank of the corresponding geometric*
 212 *design. Moreover, equality holds if and only if \mathcal{D} is isomorphic to the geometric design.*

213 Hamada's conjecture has been proved in the following cases: Hamada and Ohmori [9]
 214 established the conjecture for the design of hyperplanes in a binary projective or affine space
 215 ($q = 2, d = n - 1$). Doyen et al. [6] proved the conjecture for the design of lines in a binary
 216 projective space ($q = 2, d = 1$), as well as for the design of lines in a ternary affine space
 217 ($q = 3, d = 1$). Teirlinck [28] proved the conjecture for the design of planes in a binary affine
 218 space ($q = 2, d = 2$). Tonchev [30] proved a modified version of Hamada's conjecture using
 219 generalized incidence matrices with entries over $GF(q)$ instead of $(0, 1)$ -incidence matrices,
 220 for the classical designs having as blocks the complements of hyperplanes in $PG(d, q)$ or
 221 $AG(d, q)$ ($d = n - 1, q$ an arbitrary prime power).

222 Nevertheless, the strong version of Hamada's conjecture is not true in general: there are
 223 designs with the same parameters and the same p -rank as a classical geometric design \mathcal{D} ,
 224 but not isomorphic to \mathcal{D} . The smallest examples for this phenomenon are the quasi-sym-
 225 metric designs with the parameters of $PG_2(4, 2)$, namely, 2-(31, 7, 7) [29], which were

⁴ This is not contained in the published paper [21], but was mentioned by Mavron and McDonough to the second author when he was visiting The University of Wales at Aberystwyth.

226 already discussed in the previous section. We note that one of these 2 - $(31, 7, 7)$ designs,
 227 namely, the design supported by the minimum weight vectors in the quadratic-residue code
 228 of length 31 , was mentioned in the paper by Goethals and Delsarte [7]. The extensions of the
 229 quasi-symmetric 2 - $(31, 7, 7)$ designs are 3 - $(32, 8, 7)$ designs having the same parameters
 230 and block intersection numbers as $AG_3(5, 2)$ [29]. All these designs have the same 2 -rank,
 231 namely 16 .

232 The only other previously known parameter set for which a non-geometric design exists
 233 that has the same p -rank as the corresponding geometric design is 2 - $(64, 16, 5)$: in [10],
 234 Harada et al. found two affine 2 - $(64, 16, 5)$ designs having the same 2 -rank (equal to 16) as
 235 the classical geometric design of the planes in $AG(3, 4)$. The two exceptional designs were
 236 found as minimum weight vectors in binary codes spanned by incidence matrices of sym-
 237 metric $(4, 4)$ -nets. Mavron et al. [22] showed that one of the pseudo-geometric 2 - $(64, 16, 5)$
 238 designs from [10] can be obtained also by using a certain line spread in $PG(5, 2)$.

239 However, the weak version of Conjecture 4.1, that is, the statement that the p -rank of any
 240 design with the same parameters as a geometric design $PG_d(n, q)$ or $AG_d(n, q)$ is at least as
 241 large as that of the corresponding geometric design, is still open in general, with the exception
 242 of the few proven cases mentioned above.

243 Thus, it is rather interesting that the designs described in Theorem 2.4 in the case when q
 244 is a prime number provide the first infinite family of counterexamples to the strong version
 245 of Hamada's conjecture:

246 **Theorem 4.2** *If $q = p$ is a prime number, the pseudo-geometric designs described in The-*
 247 *orem 2.4 have the same p -rank as the geometric design $PG_k(2k, p)$.*

248 We will need two lemmas for the proof of Theorem 4.2.

249 **Lemma 4.3** *Let α be a polarity in $PG(2k - 1, q)$, where $q = p^s$ and p is a prime. The*
 250 *p -rank $r_p(\alpha)$ of the incidence matrix of the design $\alpha(\mathcal{D})$ from Theorem 2.4 satisfies the*
 251 *inequalities*

$$252 \quad r_p(\mathcal{D}) \leq r_p(\alpha) \leq \frac{1}{2} \left(\frac{q^{2k+1} - 1}{q - 1} + 1 \right), \quad (6)$$

253 where $r_p(\mathcal{D})$ is the p -rank of the geometric design $\mathcal{D} = PG_k(2k, q)$.

254 *Proof* By the construction described in Sect. 2, the design $\alpha(\mathcal{D})$ has an incidence matrix of
 255 the form

$$256 \quad M = \left(\begin{array}{c|c} M_1 & M_2 \\ \hline 0 & M_3 \end{array} \right),$$

257 where M_1 is a point by block incidence matrix of the geometric design $PG_k(2k - 1, q)$, and
 258 M_3 is a point by block incidence matrix of the geometric design $AG_k(2k, q)$. Thus, we have

$$259 \quad r_p(M_1) + r_p(M_3) \leq r_p(\alpha).$$

260 On the other hand, it follows from [1, Corollary 5.7.3, p. 186], that

$$261 \quad r_p(PG_k(2k, q)) = r_p(PG_k(2k - 1, q)) + r_p(AG_k(2k, q)).$$

262 Hence, we have

$$263 \quad r_p(\mathcal{D}) = r_p(M_1) + r_p(M_3).$$

264 This proves the left-hand side inequality in (6). To prove the right-hand side inequality in
 265 (6), we consider the complementary design $\alpha(\mathcal{D})$. By Lemma 2.2, the design $\alpha(\mathcal{D})$ has the

266 same intersection numbers as $\mathcal{D} = PG_k(2k, q)$, that is, $(q^i - 1)/(q - 1)$ for i in the range
 267 $1 \leq i \leq k$. Consequently, the block intersection numbers of the complementary design $\overline{\alpha(\mathcal{D})}$
 268 are

$$269 \quad \frac{q^i(q^{2k+1-i} - 2q^{k+1-i} + 1)}{q - 1}, \quad 1 \leq i \leq k.$$

270 Note that all these numbers are divisible by q , and that the blocks of $\overline{\alpha(\mathcal{D})}$ are of size

$$271 \quad \frac{q^{k+1}(q^k - 1)}{q - 1},$$

272 which is also divisible by q . Thus, the incidence vectors of the blocks of $\overline{\alpha(\mathcal{D})}$ span a linear
 273 self-orthogonal code of length $(q^{2k+1} - 1)/(q - 1)$ over $GF(p)$. Hence, the p -rank of the
 274 incidence matrix $(J - M)$ of $\overline{\alpha(\mathcal{D})}$, where J denotes the all-one matrix of appropriate size,
 275 does not exceed $(\frac{q^{2k+1}-1}{q-1} - 1)/2$ (note that the number of points of $\alpha(\mathcal{D})$, $(q^{2k+1} - 1)/(q - 1)$
 276 is an odd number). The columns of $J - M$ have 0 and 1 entries, and the number of 1's in each
 277 column is a multiple of p . Therefore, each column of $J - M$ is orthogonal (over $GF(p)$) to
 278 the all-one column \mathbf{j} , and consequently, the whole column space is orthogonal to \mathbf{j} . Since \mathbf{j}
 279 is not orthogonal to itself, \mathbf{j} is not in the column space of $J - M$. On the other hand, \mathbf{j} is a
 280 nonzero multiple of the sum of columns of M over $GF(p)$. This implies

$$281 \quad r_p(M) = r_p(J - M) + 1,$$

282 and therefore

$$283 \quad r_p(M) \leq \frac{1}{2} \left(\frac{q^{2k+1} - 1}{q - 1} - 1 \right) + 1 = \frac{1}{2} \left(\frac{q^{2k+1} + 1}{q - 1} + 1 \right).$$

284 This proves the right-hand side inequality in (6). \square

285 A summation formula for the p -rank of the incidence matrix of a geometric design
 286 $PG_r(n, q)$, $1 \leq r \leq n - 1$, $q = p^t$, p a prime, was found by Hamada [8]. If $r \neq 1$, $n - 1$,
 287 Hamada's formula involves some parameters that have to be computed. A simplified formula
 288 for the case when $q = p$ is a prime was found by Hirschfeld and Shaw [13, Theorem 5.10].
 289 In particular, the p -rank of $\mathcal{D} = PG_k(2k, p)$ is given by:

$$290 \quad r_p(\mathcal{D}) = \frac{p^{2k+1} - 1}{p - 1} - \sum_{i=0}^{k-1} (-1)^i \binom{(k-i)(p-1)-1}{i} \binom{k+(k-i)p}{2k-i}. \quad (7)$$

291 If $p = 2$, the linear code spanned by the blocks of $\mathcal{D} = PG_k(2k, 2)$ is a punctured
 292 Reed-Muller code of length $v = 2^{2k+1} - 1$ and order k [1, Proposition 5.3.2], so we have an
 293 alternative formula for $r_2(\mathcal{D})$ which can be written in a simple closed form, namely

$$294 \quad r_2(\mathcal{D}) = \sum_{i=0}^k \binom{2k+1}{i} = 2^{2k}.$$

295 Note that $2^{2k} = (v + 1)/2$, so the inequalities in (6) are replaced by equalities:

$$296 \quad r_2(\mathcal{D}) = r_2(\alpha) = 2^{2k} = (v + 1)/2.$$

297 Thus, the pseudo-geometric designs from Sect. 2 for $q = p = 2$ are counter-examples to the
 298 "only if" part of Hamada's conjecture.

299 In addition, the two formulas for $r_2(\mathcal{D})$ imply the following identity:

$$300 \quad 2^{2k} - 1 = \sum_{i=0}^{k-1} (-1)^i \binom{k-i-1}{i} \binom{3k-2i}{2k-i}. \quad (8)$$

301 It turns out that a similar closed formula for $r_p(\mathcal{D})$ holds for any prime number p .

302 **Lemma 4.4** *If p is any prime, the p -rank of $\mathcal{D} = PG_k(2k, p)$ is equal to*

$$303 \quad r_p(\mathcal{D}) = \frac{1}{2} \left(\frac{p^{2k+1} - 1}{p - 1} + 1 \right). \quad (9)$$

304 *Proof* We will use the following result by Hirschfeld and Shaw [13, Corollary 5.5]: if p is
 305 a prime and $C^*(k, n, p)$ is the dual of the linear code over $GF(p)$ spanned by the incidence
 306 vectors of the k -dimensional subspaces of $PG(n, p)$, $1 \leq k \leq n - 1$, then

$$307 \quad \dim C^*(k, n, p) + \dim C^*(n - k, n, p) = \frac{p^{n+1} - 1}{p - 1} - 1. \quad (10)$$

308 In the special case $n = 2k$, (10) implies that

$$309 \quad \dim C^*(k, 2k, p) = \frac{1}{2} \left(\frac{p^{2k+1} - 1}{p - 1} - 1 \right).$$

310 Note that $C^*(k, 2k, p)$ is the code having the incidence matrix of $\mathcal{D} = PG_k(2k, p)$ as a parity
 311 check matrix, hence

$$312 \quad r_p(\mathcal{D}) = \frac{p^{2k+1} - 1}{p - 1} - \dim C^*(k, 2k, p) = \frac{1}{2} \left(\frac{p^{2k+1} - 1}{p - 1} + 1 \right).$$

313 □

314 Now Theorem 4.2 follows from Lemmas 4.3 and 4.4.

315 We note that comparing (7) and (9) gives the following identity, which generalizes (8):

Corollary 4.5

$$316 \quad \frac{1}{2} \left(\frac{p^{2k+1} - 1}{p - 1} - 1 \right) = \sum_{i=0}^{k-1} (-1)^i \binom{(k-i)(p-1)-1}{i} \binom{k+(k-i)p}{2k-i}. \quad (11)$$

317 It was pointed out to us by one of the reviewers, that Eq. 11 is actually true for all positive
 318 integers p and not just for primes; it follows from a formula of J.L.W.V. Jensen [14, Eq. 18],
 319 which is given a modern setting in [20, Sect. 14.1]. Of course, with (11) in hand, Lemma 4.4
 320 is an immediate consequence of (7).

321 We finally remark that Theorem 4.2 does not extend to arbitrary prime powers q : the
 322 classical design $PG_2(4, 4)$ has 2-rank 146, whereas the pseudo-geometric design obtained
 323 from a polarity has 2-rank 154.

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